

Sweap: Reactive Synthesis for Infinite-State Integer Problems

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∞ -state Synthesis: Sweap's Setting

Novel features are *highlighted*

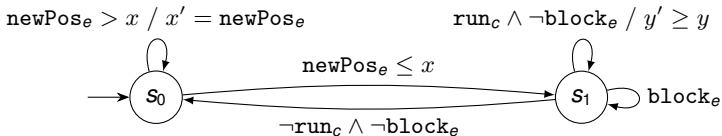
Arena: $A = \langle V, \mathbb{E}, \mathbb{C}, val_0, \delta \rangle$ Game: $\langle A, \varphi \rangle$

- V : state variables (integers)
- $\mathbb{E}$: inputs (Boolean **or integers**), $\mathbb{C}$: outputs (Boolean)
- val_0 : initial valuation of state variables (**optional**)
- δ : transition relation (**nondeterminism** resolved by controller)
- $\varphi \in \text{LTL}(\mathbb{E} \cup \mathbb{C} \cup \mathcal{Pr})$, $\mathcal{Pr}$ = set of predicates over V
- E.g., $G ((x = 0) \Rightarrow F (x = 5))$

Realizability modulo arena

- Let environment $\mathcal{E}$ set the values of predicates.
- A trace is concretisable iff $\mathcal{E}$'s values for predicates agree with the valuation of V at every step.
- φ is realizable modulo A iff there is a strategy that preserves φ on every concretisable trace

Example: Infinite Race



$$V = \{x \stackrel{\text{init}}{=} 0, y \stackrel{\text{init}}{=} 0\}; \quad \mathbb{E} = \{\text{newPos}_e, \text{block}_e\}; \quad \mathbb{C} = \{\text{run}_c\}$$

$$\textbf{Goal: } GF(s_1 \wedge \neg \text{block}) \implies GF(s_0 \wedge (y > x))$$

Assumption

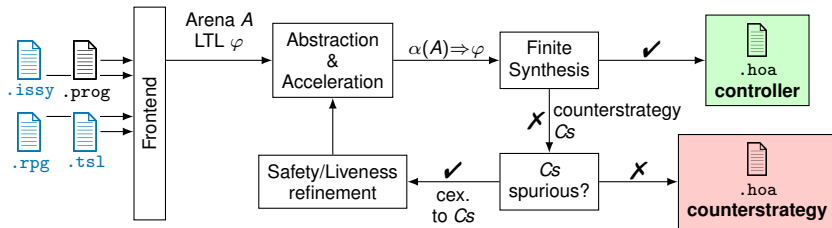
Environment cannot stay in s_0 forever,
cannot block controller forever

Guarantee

Controller cannot stay in s_1 forever,
must be ahead of the environment
 ∞ often

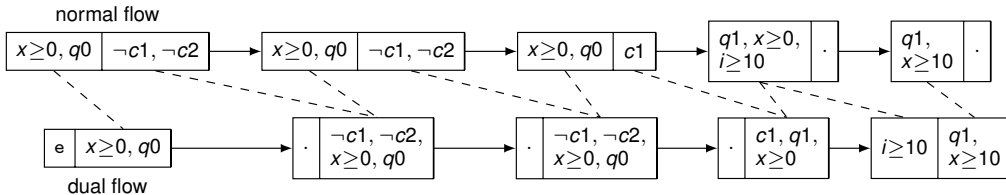
Sweap's CEGAR Approach

Novel features are *highlighted*

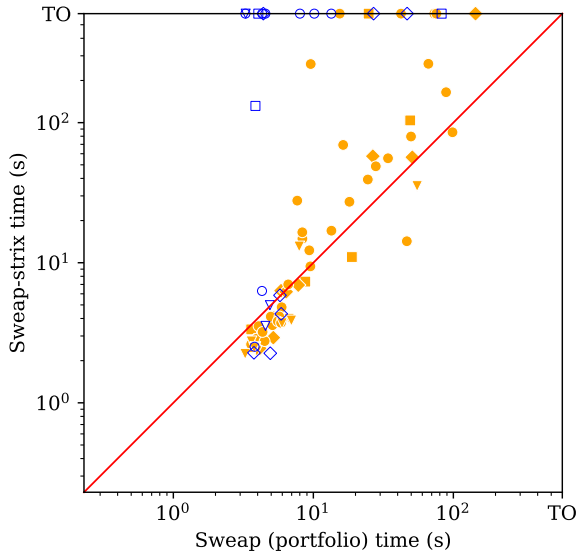


- Liveness refinement prunes unconcretisable infinite loops
- Acceleration: fairness assumptions on truth values of predicates
- Finite synthesis backends: Strix, [SemML](#) (new default)
- Main algorithmic contribution: [dualization](#)

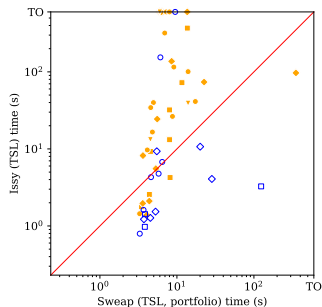
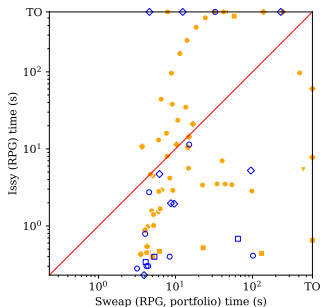
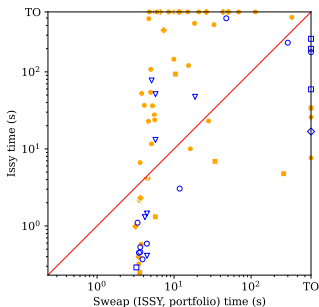
- Recall: we let $\mathcal{E}$ choose predicate values
- This approach does **not terminate** on some problems
- Specifically, those that do not admit **finite counterstrategies**
- Dualization **flips** control of the predicates to the **controller**
- However, this slightly breaks the order of play in the game
 $\Rightarrow$ We also turn φ into equirealisable φ^{dual}
 (add X s, put additional constraints on $\mathcal{E}$)



- Sweap-strix
 - Uses Strix, no dualization
(Roughly same as CAV'25 + bug fixes and optimisations)
- Sweap (portfolio)
 - Uses SemML
Best of normal + dual workflows
- 95 benchmarks (72 R, 23 U) in Sweap's native format
- 19 new verdicts (12 by dualization)
- Strix faster on smaller instances



Format	# Benchmarks	Tool	Realisable			Unrealisable		
			S	F	U	S	F	U
ISSY	115 (89 R, 26 U)	Issy Sweep	42	18	2	20	14	3
			70	53	40	21	5	4
RPG	86 (66 R, 20 U)	Issy Sweep	52	34	2	16	14	0
			57	25	7	20	6	4
TSL	65 (45 R, 20 U)	Issy Sweep	32	17	2	14	10	0
			38	21	8	17	7	3



- CEGAR promising for ∞ -state synthesis
- Dualization helps on unrealizable specifications
- Issy still (sometimes) faster, less memory hungry

Future work:

- Go beyond black-box finite synthesis
- Go beyond LIA
- Go beyond “monolithic” game arena

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<https://github.com/shaunazzopardi/sweap/>